

METHOD AND APPARATUS FOR ADAPTIVE SIGNALING IN A QAM COMMUNICATIONS SYSTEM

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FIELD OF THE INVENTION

The present invention relates generally to communications systems, and more particularly to providing signaling through puncturing data symbols in a quadrature amplitude modulation communication system.

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BACKGROUND OF THE INVENTION

Various wireless communication systems are known in the art. In multiple access wireless communications systems, such as code division multiple access (CDMA), a base station transmits multiple signals to individual mobile stations. The base station transmits multiple signals on a forward link that typically includes separate data and signaling channels. Similarly, the mobile station transmits data and signaling via a reverse link to the base station.

Signaling channels used in the communications system are often used for tasks such as power control and for sharing system information such as data frame structures. The signaling channels transmit at a low rate and require minimal latency and high reliability since the information transmitted via these channels is used to control the communications system. High latency or erroneous data in the signaling channels may cause the communications system to become unstable, which severely degrades the system capacity. This degradation may be further exacerbated in a multiple access communication system that achieves high data rates through use of quadrature amplitude modulation (QAM) along with Turbo encoders and decoders.

As mentioned previously, signaling channels may be implemented using a separate dedicated channel, which has its own convolutional encoders and decoders. Separate dedicated channels, however, require system resources that may be costly and not readily available and diminish resources available to other channels. Additionally, 5 the separate dedicated channels incur a convolutional decoder delay. In order to reduce signaling overhead and avoid signaling delay, another method of implementing a signaling channel is to puncture bits on top of existing high data rate data channels. However, signaling accomplished through puncturing on top of existing data channels can degrade the performance of the high data rate channel and can also suffer a loss of 10 reliability because there is no coding gain in this method. In this instance, attempts to increase reliability have included repetition of signaling bits, which consumes even greater amounts of system resources.

Therefore, a need exists for signaling channels in a high data rate QAM communication system that have high reliability and low decode delay while utilizing 15 minimal system resources.

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 is an illustration of an exemplary wireless communication system in accordance with an embodiment of the present invention.

20 FIG. 2 is a block diagram of communications system architecture in accordance with an embodiment of the present invention.

FIG. 3 is an illustration of a Gray-coded Karnaugh mapped QAM constellation utilized by the QAM mapper of FIG. 2.

FIG. 4 is a table showing the correlation of log-likelihood ratio gain for each bit position in a data symbol based on values of other bits within the data symbol in accordance with an embodiment of the present invention.

FIG. 5 is an illustration of various puncturing schemes utilized for puncturing 5 particular bits in a data symbol in accordance with embodiments of the present invention.

DETAILED DESCRIPTION OF THE INVENTION

To address the need for signaling channels in a high data rate QAM 10 communication system that have high reliability and low decode delay while utilizing minimal system resources, a method and apparatus is provided that uses quadrature amplitude modulation (QAM) with bit-wise decoders where a log-likelihood ratio (LLR) is computed for each bit at a receiver portion. In particular, a transmitter in a communication system punctures bits of a data symbol that includes multiple data bits 15 prior to modulation of the data symbol by a QAM mapper. The QAM mapper may employ a Gray-coded Karnaugh mapped constellation. In addition, the transmitter employs adaptive puncturing wherein specific bits within a data symbol are selectively punctured to achieve higher LLR gains that result from characteristics of the Karnaugh mapping scheme. The increased LLR gain affords higher reliability for signaling 20 without consuming additional system resources.

The present invention can be better understood with reference to FIGs. 1-5.

FIG. 1 illustrates an exemplary communications system 100 in accordance with an embodiment of the present invention. Communications system 100 includes a base

station 102 that is capable of transmitting multiple signals to each of multiple mobile subscriber units (MSs) 104 (one shown), such as cellular telephones, radiotelephones, or wireless data modems. Base station 102 transmits multiple signals on a forward link that includes data and signaling channels. Similarly, MS 104 transmits data and
5 signaling via a reverse link to the base station 102.

FIG. 2 is a block diagram of an architecture of a communications system 200, such as communications system 100, that implements the signaling channel by adaptively puncturing signaling bits into a data symbol. Communications system 200 includes a transmitter (indicated by bracket), such as base station 102. Communications
10 system 200 further includes a receiver (indicated by bracket), such as MS 104, on the other side of a transmission interface 208 that receives transmitted data symbols and that extracts the signaling bits from the punctured data symbols after the computation of a log-likelihood ratio (LLR) for each bit of a received data symbol. However, those who are of ordinary skill in the art realize that base station 102 and MS 104 are each capable
15 of operating as either a transmitter or a receiver with respect to the embodiments of the present invention.

On the transmitter side of communications system 200, a stream of data, preferably in a binary format such as bits, is input to an encoder 202. Encoder 202 encodes the data stream pursuant to a prescribed data coding scheme and routes the
20 encoded data to a QAM mapper 204. The type of encoder is not critical to, nor is specific encoding necessary for, the present invention. In another embodiment of the present invention, the data stream may be entered directly to QAM mapper 204. In either case, prior to inputting the data stream to QAM mapper 204, signaling data (i.e.,

signaling bits) is punctured into the data stream by a signal bit encoder 206 according to a predetermined puncturing scheme, which puncturing scheme is described in greater detail below.

QAM mapper 204 maps the input data stream to points in a multi-dimensional constellation. In order to map the data stream, QAM mapper 204 groups the input data stream into multiple groups of P bits, that is, into multiple P-tuples wherein each P-tuple may be thought of as a data symbol. QAM mapper 204 modulates each of the multiple P-tuples, or data symbols, by mapping the P-tuple to a corresponding point out of M possible points in a predetermined M-ary QAM constellation, wherein $M = 2^P$.

To this end, the predetermined QAM constellation that includes the M possible points is defined within a multi-dimensional space, preferably a complex two-dimensional (I/Q) space. Each point within the two-dimensional space may be thought of as a vector sum of two scaled basis vectors. The two scaled basis vectors respectively correspond to an in-phase (I), or a real (r), component and a quadrature (Q), or an imaginary (i), component of the constellation point, or corresponding data symbol. The respective amplitudes of the two basis vectors used to define a particular point may be thought of as two-dimensional coordinates of the point.

Preferably the mapping scheme employed by QAM mapper 204 includes a Gray-coded Karnaugh mapping scheme, which is utilized to effect the predetermined puncturing scheme of the present application and simplifies hardware mapping and slicing functions. However, those who are of ordinary skill in the art realize that other QAM mapping schemes may be used herein without departing from the spirit and scope of the present invention. After modulation of each data symbol by QAM mapper 204, the modulated data symbols are transmitted via a communication interface 208 to the

receiver. Communication interface 208 is preferably wireless radio frequency (RF) transmission interface, but could also be any other transmission interface means such as land lines.

At the receiver of communications system 200, a Symbol-to-LLR calculator 210 receives each modulated data symbol and demodulates the modulated data symbol based on a determined log-likelihood ration (LLR) for each bit of the multiple bits corresponding to the modulated data symbol. Symbol-to-LLR calculator 210 determines an LLR for each bit based on a predetermined algorithm that is discussed in greater detail below. The demodulated data symbol is then output from Symbol-to-LLR calculator 210 to both a zero-fill block 214 and a signaling channel decoder 212. When a data symbol has been punctured, zero-fill block 214 fills punctured bit locations in the data symbol with a soft value of zero corresponding to an equal likelihood of a bit value of a "0" or a "1". Each demodulated data symbol is then decoded by a bit-wise decoder 216 to recover the data bits encoded by encoder 202.

FIG. 3 is an illustration of an exemplary QAM constellation 300 in accordance with an embodiment of the present invention. QAM constellation 300 preferably comprises a Gray-coded Karnaugh constellation; however, mapping schemes other than Karnaugh mapping may also be used herein without departing from the spirit and scope of the present invention. QAM constellation 300 is a 64-QAM square constellation (i.e., $M = 64$) wherein each of the 64 constellation points corresponds to a particular P-tuple, or data symbol, of 6 bits (i.e., $P = 6$). However, those of ordinary skill in the art realize that any square Karnaugh mapped constellation for an M-ary QAM system may be used herein without departing from the spirit and scope of the present invention.

Each P-tuple, or 6-bit sequence, is of the form “ $i_1 q_1 i_2 q_2 i_3 q_3$ ”, labeled from the most significant bit to least significant bit. Each “ i ” bit is independent of the imaginary, or quadrature (Q), axis 304 and each “ q ” bit is independent of the real, or in-phase (I), axis 302.

5 For example, for any given point in QAM constellation 300, movement to other points in a direction that is parallel to I-axis 302 produces a change in one or more “ i ” bits in the corresponding 6-bit sequence but does not produce a change in any of the “ q ” bits. As another example of this independence, for each point included in the first column of points to the right of Q-axis 304, bit i_1 in each corresponding data symbol is
10 “0”, bit i_2 in each corresponding data symbol is “0”, and bit i_3 in each corresponding data symbol is “1”. As a further illustration of this property of the Karnaugh mapped QAM data symbols, the lines shown at the top of, and to the left of, constellation 300 indicate subsets of constellation points where an indicated bit of a corresponding data symbol is a “1”, and an absence of a line indicates of a subset of constellation points
15 where an indicated bit of a corresponding data symbol is a “0”. For example, the subset of points in constellation 300 wherein the bit i_1 of a corresponding data symbol is a “1” lies in the left hand side of the I/Q plane whereas the subset of points wherein the bit i_1 of a corresponding data symbol is a “0” lies in the right hand side of the I/Q plane.

20 Symbol-to-LLR calculator 210 determines an LLR for each of the bits within a data symbol to determine whether the bit is more likely to be a “0” or a “1”. For a received data symbol y received at time k (i.e., y_k), the LLR of a j -th bit of the received symbol is determined by the following relationship:

$$LLR(u_{k,j}) = \log \left\{ \frac{P(u_{k,j} = 0 | y_k)}{P(u_{k,j} = 1 | y_k)} \right\} \quad (1)$$

where “ $u_{k,j}$ ” corresponds to a hypothesized j -th bit of the transmitted data symbol based on the received symbol y_k , “ $P(u_{k,j} = 0 | y_k)$ ” corresponds to a probability that the hypothesized j -th bit is a value “0” given the received data symbol y_k , and “ $P(u_{k,j} = 1 | y_k)$ ” corresponds to a probability that the hypothesized j -th bit is a value “1” given the received data symbol y_k .

Equation (1) reduces to the expression:

$$LLR(u_{k,j}) = \log \left\{ \frac{\sum_{u_k: u_{k,j}=0} P(y_k | u_k)}{\sum_{u_k: u_{k,j}=1} P(y_k | u_k)} \right\} \quad (2)$$

where it is assumed that all transmitted data symbols are equiprobable. The probability “P” can then be represented by the following relationship:

$$P(y_k | u_k) = p(y_k^r, y_k^i | u_k^r, u_k^i) = p(y_k^r | u_k^r)p(y_k^i | u_k^i) \quad (3)$$

where “r” corresponds to the real component, and “i” corresponds to the imaginary component, of each of a data symbol u_k , selected as described below, and the received data symbol y_k , and where perfect channel correction is assumed such that the real component is independent of the imaginary component.

Assuming additive Gaussian noise, the equation (3) can be rewritten as the following expression:

$$P(y_k | u_k) = \frac{1}{2\pi\sigma^2} \exp \frac{-D_k^2}{2\sigma^2} \quad (4)$$

where the parameter D_k^2 is the squared Euclidean distance in the complex I/Q plane between a point corresponding to the received data symbol y_k and a data symbol, that is,

u_k , corresponding to one of the M points in the QAM constellation, such as constellation 300, selected as described below. It should be noted that the point in the complex I/Q plane corresponding to received data symbol y_k is unlikely to be one of the M points in the QAM constellation. The parameter σ^2 is the variance of the Gaussian noise.

5 With respect to each bit of the received data symbol y_k , two points and corresponding data symbols u_k are selected as follows from the QAM constellation of M points. A first point corresponds to the data symbol u_k whose j -th bit, that is, u_{kj} , is a value “0”, and which point, out of all of the constellation points whose j -th bit is a “0”, is nearest to the point corresponding to y_k . A second point corresponds to the data 10 symbol u_k whose j -th bit, *that is*, u_{kj} , is a value “1”, and which point, out of all of the constellation points whose j -th bit is a “1”, is nearest to the point corresponding to y_k .

For each constellation point and corresponding data symbol selected as described above, a squared Euclidean distance D_k^2 between the point and a point corresponding to the received data symbol y_k can be calculated by summing the squares of the differences 15 between the real (r) components, and between the imaginary (i) components, of the points, as represented by the following relationship:

$$D_k^2 = |y_k - u_k|^2 = (y_k^r - u_k^r)^2 + (y_k^i - u_k^i)^2 \quad (5)$$

Substituting equation (4) into equation (2) allows calculation of the LLR based on the squared Euclidean distances and yields the following expression:

$$20 \quad LLR(u_{k,j}) = \log \sum_{u_k:u_{k,j}=0} e^{\frac{-D_k^2}{2\sigma^2}} - \log \sum_{u_k:u_{k,j}=1} e^{\frac{-D_k^2}{2\sigma^2}} \quad (6)$$

Equation (6) may then be approximated by simply taking the difference between the minimum squared distances between the received data symbol y_k and each of the two

selected data symbols u_k , having respective j -th bit values of “1” and “0”. That is,

equation (6) may be reduced to the following expression to approximate the LLR:

$$LLR(u_{k,j}) = \min_{u_k: u_{k,j}=1} [D_k^2] - \min_{u_k: u_{k,j}=0} [D_k^2] \quad (7)$$

By employing equation (7), the receiver of communication system 200 can

5 determine an LLR with respect to the j -th bit of the received data symbol. Based on the LLR, the receiver can then determine a value of the j -th bit of the transmitted data symbol, thereby recovering the transmitted data symbol based on the received data symbol. The process of determining an LLR for the j -th bit of the received data symbol may be summarized in the following steps:

10 a) determining a first minimum squared Euclidean distance between the point corresponding to the received data symbol y_k and a nearest constellation point whose j -th bit, that is, $u_{k,j}$, is a value “1”,

15 b) determining a second minimum squared Euclidean distance between the point corresponding to the received data symbol y_k and a nearest constellation point whose j -th bit, that is, $u_{k,j}$, is a “0”, and

20 c) determining a difference between the first minimum squared Euclidean distance and the second minimum squared Euclidean distance.

If the difference is positive, then the j -th bit of the transmitted data symbol is

most likely a binary value of “0”. If the difference is negative, then the j -th bit is most

likely a binary value of “1”. These computed LLR values can, in turn, be fed into a convolutional decoder or a Turbo decoder, for example, that then determines the most likely sequence of bits. As an alternative, a simple slicer may be used to decode each bit by simply taking the sign of the LLR.

Several properties of the particular LLRs computed arise from using a Karnaugh-mapped construction of a QAM constellation, such as QAM constellation 300. For example, when a 6-bit data symbol is being mapped to a point in the QAM constellation, the LLRs of the even bits (i.e., q_1 , q_2 , and q_3) are independent from the 5 LLRs of the odd bits (i.e., i_1 , i_2 , or i_3). Hence, whether the value of an odd bit "i" is a "1" or "0", it has no effect on the LLR of an even bit "q", and vice versa. Furthermore, the odd bits (i.e., i_1 , i_2 , or i_3) normally determine a position in the I-axis direction, which has no effect on the LLR for any even bit q_1 , q_2 , or q_3 since all of the even bits "q" at a particular location in the Q-axis direction will have the same bit 10 value irrespective of where the data symbol lies along the I-axis. It is noted, however, the value of the odd bit or even bit can effect the LLR for the other odd or even bits, respectively.

Another feature that arises from utilizing a Karnaugh-mapped QAM constellation is that the two most significant bits (i.e., i_1 and q_1) have a higher average 15 LLR than the lesser significant bits (i.e., i_2 and q_2). These middle bits, in turn, have a higher average LLR than the two least significant bits (i.e., i_3 and q_3). The reason for the differing average LLRs from the most significant to the least significant bits is due to the fact that the most significant bits i_1 and q_1 have the largest continuous coverage in the Karnaugh map (e.g., four columns or rows of points, or corresponding data 20 symbols, as are seen in FIG. 3), whereas lesser significant bits i_2 and q_2 for example, have lesser coverage (e.g., two rows or columns of points or corresponding data symbols). The least significant bits i_3 and q_3 have the least coverage in the Karnaugh map given a data symbol having particular values of the most significant and lesser

significant bits, wherein the coverage is only one row or column of points or corresponding data symbols. Thus, the least significant bits i_3 and q_3 will always have an LLR associated with the distance of one point or corresponding data symbol.

Similarly the lesser significant bits i_2 and q_2 will have LLRs associated with distances of

5 one or two data symbols and the most significant bits i_1 and q_1 will have LLRs associated with distances of 1, 2, 3 or 4 points or corresponding data symbols.

The above-described properties of a Karnaugh mapped QAM constellation can be utilized to implement an adaptive puncturing scheme to transmit signaling information via a punctured data symbol in a simple and efficient manner. That is, bits selected to be punctured with signaling information may be chosen to yield the highest LLR or, alternatively, a minimally acceptable LLR gain that will effect a requisite level of reliability for transmitting the signaling information. The greater the value of the LLR, the greater a level of noise that is required to cause a bit error. Hence, higher LLRs are more immune to noise introduced during transmission of the data symbols.

15 FIG. 4 illustrates a table of LLR gains of bit positions and values for the Karnaugh mapped, Gray-coded constellation of FIG.3 assuming a noiseless environment. For purposes of clarity and assuming a 64-QAM system as an example, the 6 bits in a data symbol having the bit sequence “ $i_1 q_1 i_2 q_2 i_3 q_3$ ” are hereinafter referred to as corresponding sequential bit positions “ $b_5 b_4 b_3 b_2 b_1 b_0$ ”. Thus, for bit positions b_0 and b_1 , which are the least significant bits, the table in FIG. 4 illustrates that the values of these bits do not affect the magnitude of the LLR for these bits. However, the values of bit positions b_0 and b_1 affect the LLRs of the other bits in the data symbol. Thus, if these bits are punctured, a resultant gain in magnitude of the

LLRs of the other bits occurs . For instance, when bit b_0 is a "0" the resultant LLR for bit b_2 is a minimal value of $+/- 4 D^2$. However, when bit b_0 is a "1" the resultant magnitude for bit b_2 increases to $+/- 16 D^2$. Similarly for odd bit b_3 the LLR depends on the value of odd bit b_1 .

5 For the most significant bit positions b_4 and b_5 , various combinations of values of the least significant and lesser significant bits yield varying effects on the LLRs of the most significant bits. As may be seen in the table of FIG. 4, when both the lesser and least significant bits are of value "1" the most significant gain in the LLR is effected whereas when the lesser significant bit has a value of zero and the least

10 significant bit has a value of one (i.e., a bit combination resulting in a data symbol that corresponds to a constellation point that is closest to an axis), the LLR magnitude is at the minimal value of $+/- 4 D^2$.

Based on the known LLR values corresponding to particular combinations of bits in the data symbol, a data symbol having particular bit values may be punctured in order to achieve a desired LLR and a desired level of reliability for the signaling bits punctured within the symbol. Therefore, a data symbol may be punctured periodically at a known time such that the receiver is aware of where the punctured bits are located. For example, of the 6 bits comprising a transmitted data symbol that is modulated based on a 64-QAM system, one of the most significant bits (e.g., b_5) is punctured since the

15 most significant bits inherently have average LLRs superior to the lower four bits in the data symbol as was illustrated in the table of FIG. 4. By simply puncturing the most significant bit, a resultant average LLR magnitude of $30 D^2$ is achieved, which is superior to the average LLR magnitude of $14.67 D^2$ for all bits and corresponds to a 6.2

dB gain. FIG. 5 illustrates three possible bit sequences for a punctured data symbol where "X" indicates the punctured signaling bit that carries the signaling information (i.e., the signaling channel). In the first case, most significant odd bit b_5 is shown punctured without regard to the values of the other odd bits b_3 and b_1 . As discussed 5 previously, this results in an average LLR magnitude of $30 D^2$ (i.e., the average of the four LLR values corresponding to the four possible combinations of bits b_3 and b_1).

If a higher degree of reliability is required for the signaling channel the bit b_3 may be punctured with a binary value of "1" as indicated in the second sequence of FIG. 5, which will raise the average LLR magnitude for the signaling bit (i.e., b_5) to $50 D^2$ (i.e., the average of the LLR values $64 D^2$ and $36 D^2$ for the two possibilities where 10 bit b_1 is a "1" or a "0") . The increase in the LLR magnitude to an average of $50 D^2$ corresponds to a 10.6 dB gain. Similarly, in order to achieve the highest degree of reliability possible for the 64-QAM system illustrated, both odd bits b_3 and b_1 may be punctured with a binary value of "1" as shown t he third sequence illustrated in FIG. 5. 15 This will result in raising the average LLR magnitude to $64 D^2$ for signaling bit b_5 , which corresponds to a 12.8 dB gain.

It is noted that the bit sequence described above is an arbitrary sequence and that other bit sequences, as well as other QAM mapping schemes, may be employed herein without departing from the spirit and scope of the present invention. For example, a 20 mapping where the least significant bits in a data symbol correspond to the largest area of coverage on the mapping constellation could be used. In this case, the least significant bit positions in the data symbol could be selectively punctured with signaling information since these positions would inherently have the highest LLRs due to their

coverage of the constellation. Further, the lesser significant and most significant bits could be punctured to yield gains in the LLRs of the least significant bits correlative to the puncturing of lesser and least significant bit positions of the Karnaugh mapped system described previously. Other mappings where other lesser significant bits or 5 combinations of most, lesser or least significant bits represent the highest area of constellation coverage may also be envisioned.

Based on the foregoing, the degree of reliability for the signaling channel can be selectively set, dependent on the adaptive puncturing of particular bits in a data symbol and the number of those bits actually punctured. To achieve higher signaling reliability, 10 those bit positions having the highest LLRs given a particular mapping scheme (e.g., those bit positions in a data symbol having the greatest constellation area coverage) can be chosen for puncturing with signaling information. However, the more bits punctured results in greater degradation to the high data rate QAM channel. Nonetheless, convolutional or Turbo encoding of the high data rate information is capable of 15 protecting the information data bits against single-bit puncturing such that degradation is minimal. Additionally, the signaling bits themselves are easily decoded by signaling channel decoder 212 without incurring any Turbo or convolutional decoder delay.

In operation, the system of FIG. 2 effects a signaling channel by first determining a desired reliability level (i.e., the desired number of bits to be punctured 20 in a data symbol). This selection, however, need not be static, but may be varied continuously dependent on the level of noise present in the system. Signaling bit encoder 206 selects particular data symbols to puncture periodically based on a predetermined time or any other criterion. Irrespective of the methodology for deciding

which data symbol or symbols of multiple data symbols are punctured, information concerning which data symbols are punctured is common to both the transmitter and receiver such that the receiver knows which data symbols it receives are punctured with signaling channel information. Preferably, the actual signaling information is contained

5 in most significant bit b_5 and the selective reliability of the signaling channel is set by puncturing either bit b_3 or bits b_3 and b_1 , but may be different for other alternate mapping schemes and symbol sequences. Once a data symbol is punctured by the signaling bit encoder 206, the punctured data symbol, along with non-punctured data

10 symbols, are modulated by QAM mapper 204 according to the predetermined mapping scheme. Preferably this mapping scheme is a Gray-coded Karnaugh map as exemplified in FIG. 3. The modulated data symbols are then transmitted by an interface 208, such as a wireless interface, to a receiver.

Modulated data symbols are received by the receiver and are input to Symbol-to-LLR calculator 210. Symbol-to-LLR calculator 210 determines, for each bit of each

15 received data symbol, the minimum squared Euclidean distance between a point in the complex plane corresponding to the received data symbol and each of a constellation point corresponding to a data symbol where the value of the bit is a “1” and a constellation point corresponding to a data symbol where the value of the bit is a “0”. The difference of these minimum squared Euclidean distances yields the LLR. When a

20 punctured data symbol is output by Symbol-to-LLR calculator 210 that is known by the receiver to be punctured according to the predetermined puncturing scheme, signaling channel decoder 212 simply reads the value of the computed LLR for the punctured bit (here bit b_5). If the LLR for the punctured is positive then the signaling bit is decoded

as a "0", for example. Conversely, if the LLR is negative then the signaling bit is decoded as a "1". It is noted that this could be alternatively be the opposite of the foregoing, dependent on the particular mapping scheme utilized. The punctured symbol is then input to zero-fill punctured bits inserter 214, which replaces the punctured bits 5 within the punctured data symbol with a soft value of zero corresponding to an equal likelihood of a bit value of a "1" or a "0". The LLR values are then delivered to bit-wise decoder 216 that employs convolutional decoding, Turbo decoding or any other known decoding methodology known in the art to recover the data bits.

The puncturing scheme according to the teachings of the present invention may 10 be applied to any other M-ary square constellation where M is an integer power of 2 and is also greater than 4. As noted previously, the constellation for the QAM mapping may employ some form of Karnaugh mapping as that illustrated in FIG. 3 but need not be identical to this mapping or may also employ a mapping that is not Karnaugh mapping or Gray-coded. With other M-ary square constellation QAM systems, the bit 15 location or locations having the highest LLRs may be punctured with the signaling channel data. In addition, other bits that are known to affect the LLR of the bit locations having the highest LLRs may be punctured with particular binary values in order to obtain higher degrees of LLR gain.

The above teachings of the present invention may be utilized in any multiple- 20 access communication system employing QAM with bit-wise decoders and a low data rate channel for signaling purposes. Use of the bit puncturing scheme according to the teachings of the present invention affords reliability and reduces latency for signaling bits in high data rate channels. Adaptive puncturing according to the teachings of the

present application affords also a simple solution that may be easily implemented in hardware and affords flexibility to adjust the reliability to support the minimum quality of service required. It will be further apparent to those skilled in the art that other embodiments other than the specific disclosed embodiments described above may be devised without departing from the fair scope of the appended claims and their equivalents.